

# Can you still trust your network card?

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# Modern network cards

They do not only connect the host to the network.

- ▶ Hardware architectures are complex:
  - ▶ several processors,
  - ▶ different kind of memories,
  - ▶ multiple network interfaces;
- ▶ Embedded softwares (firmwares) do more than you think:
  - ▶ remote administration: ASF, IPMI, AMT, etc,
  - ▶ TCP segmentation offloading,
  - ▶ radio with temporal constraints: GSM, 802.11, etc.

## Impacts on security

If an attacker can execute arbitrary code on the card, she can do virtually anything:

- ▶ stop processing packets;
- ▶ drop some packets;
- ▶ ARP/DNS cache poisoning;
- ▶ implement SSLstrip-like attacks;
- ▶ attacks hosts on the LAN;
- ▶ replace the firmware;
- ▶ attack the host (read/write access to the main memory).

A "must read" on this topic:

Arrigo Triulzi, PacSec08, "Project Maux Mk.II"

# What will be described today

We will present:

- ▶ architectures of modern network cards;
- ▶ remote administration protocols used in these cards;
- ▶ an actual vulnerability that we discovered;
- ▶ tools developed to debug a *Broadcom NetXtreme* card;
- ▶ exploitation proof of concept and demo;
- ▶ mitigations and workarounds to this attack.

# What won't be described today

This is not about:

- ▶ driver bugs;
- ▶ OS vulnerabilities.

And please note that:

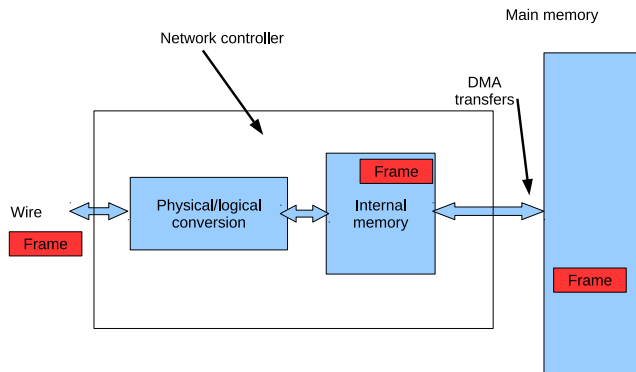
- ▶ we won't provide the network packets and the tools used in the demo;
- ▶ we worked with the vendors and they issued a patch for the vulnerability.
  - ▶ CVE-2010-0104: *HP Small Form Factor or Microtower PC with Broadcom Integrated NIC Firmware, Remote Execution of Arbitrary Code*

<http://www.ssi.gouv.fr/trustnetworkcard>

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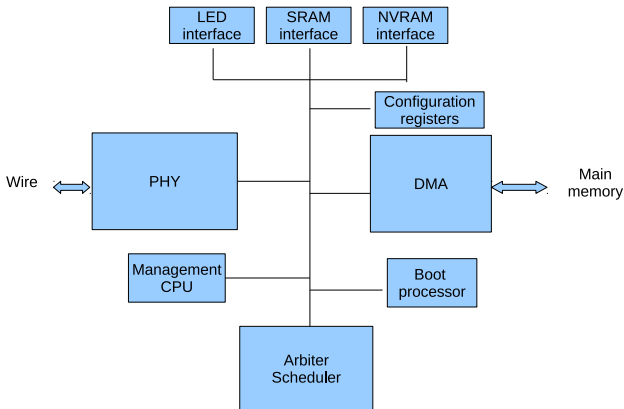
## Internal architecture of a network card

- ▶ PHY: send and receive signals on the wire;
- ▶ DMA-engine: exchange packets with the host;
- ▶ negotiation and link control (speed, duplex), etc.



## Internal architecture of NetXtreme network card

- ▶ performs various operations on packets;
- ▶ offloads work from host;
- ▶ needs hardware and software to do that;
- ▶ runs as an intercepting proxy.



## RX RISC

*An on-chip RISC processor is provided for running value-added firmware that can be used for **custom frame processing**. The on-chip RISC operates independently of all the architectural blocks; essentially, RISC is available for the **auxiliary processing of data streams**.*

- ▶ a MIPS CPU on the card;
- ▶ it has access to major components:
  - ▶ shared memory,
  - ▶ incoming and outgoing packets,
  - ▶ PCI Configuration Space,
  - ▶ SMBus;
- ▶ it executes a *firmware*.



# The firmware in the NetXtreme

Different firmwares:

- ▶ ASF (Alert Standard Format protocol);
- ▶ TSO (TCP Segmentation Offloading).

The firmware is:

- ▶ loaded from an EEPROM;
- ▶ or by the driver from the filesystem:
  - ▶ Linux driver only has TSO,
  - ▶ Windows drivers rarely have a firmware,
  - ▶ the firmware seems protected;
- ▶ loaded to memory (SRAM) during execution.

## Internal memory

- ▶ the card internal memory is mapped to the host one;
- ▶ internal memory can be accessed through a 32kb window;
- ▶ this window can be moved to read the whole internal memory space from the host.

# Alert Standard Format (ASF) 1.0 1/3

## ASF

- ▶ transmits alerts/events using the network:
  - ▶ hard disks failures, BIOS errors, ...
  - ▶ heartbeats ("machine is up");
- ▶ must operate if everything else fails (dead hard disks, OS).

The network card receives events from others devices using the SMBus (*System Management Bus*).

## RMCP

- ▶ ASF uses a protocol called *Remote Management and Control Protocol*;
- ▶ RMCP can query system state;
- ▶ RMCP allows to remotely start, stop or reboot computers.

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## Alert Standard Format (ASF) 1.0 (2/3)

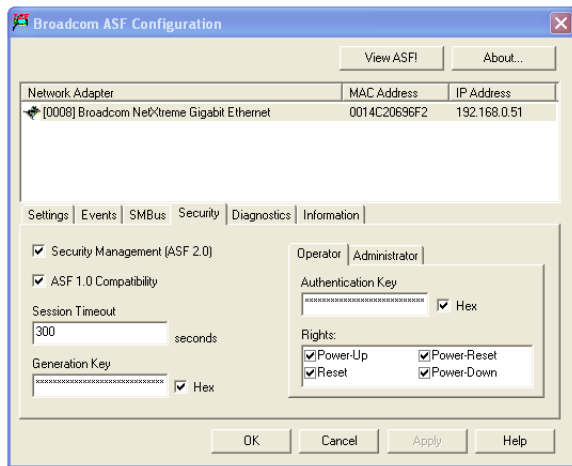
- ▶ firmware parameters must be configured: IP address, netmask, heartbeats frequency,
  - ▶ a tool is provided by network cards vendors;
- ▶ a specific ACPI table is used by ASF;
- ▶ ASF can be deactivated from the BIOS on some hardware;
- ▶ at least one boot on an ACPI-enabled OS is mandatory.

## ASF 1.0 (3/3)

### Security

- ▶ no security interfaces;
- ▶ vendors are discouraged to implement their own, proprietary security interface;
- ▶ security issues should be addressed at network infrastructure level.

# ASF configuration



- ▶ IP addresses;
- ▶ RMCP;
- ▶ permissions;

# Alert Standard Format (ASF) 2.0

ASF 2.0 adds a new protocol: RSP

- ▶ RMCP security-extensions protocol;
- ▶ adds authentication and integrity protection;
- ▶ no encryption.

This presentation focuses on ASF 2.0.

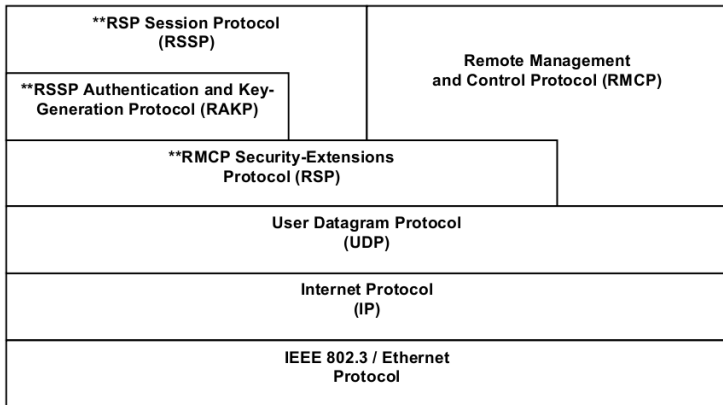
## RMCP in ASF 2.0

- ▶ messages are sent over UDP;
- ▶ traffic must be either sent on:
  - ▶ the legacy port 623/udp: no authentication, no integrity,
  - ▶ the secure port 664/udp: RMCP messages are carried inside *RMCP Security-Extensions Protocol (RSP)*;
- ▶ the network card grabs traffic on these ports, analyzes RMCP packets, and replies to queries.

The network card must implement the following stack:  
IP/UDP/RSP/RMCP.



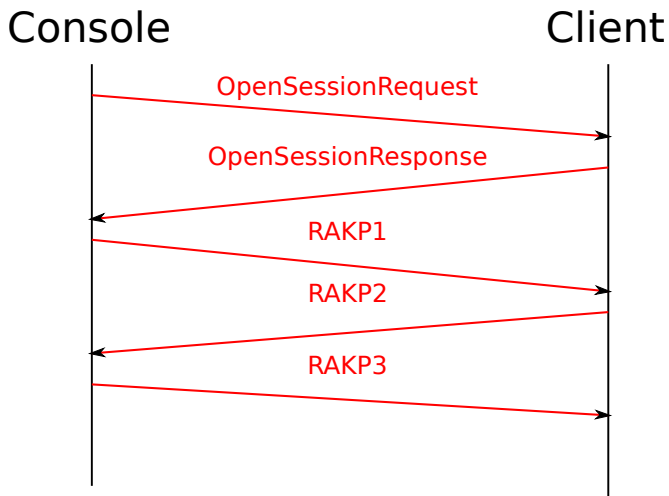
# Protocols stacks



## RMCP Security-Extensions Protocol (RSP)

- ▶ RSP adds mutual authentication of the remote console and the client;
  - ▶ the **console** is the management device used by the administrator,
  - ▶ the **client** is the remote workstation;
- ▶ messages are authenticated using pre-shared HMAC-SHA1 keys;
- ▶ in order to send a message on the secure port, the console must:
  - ▶ open a session,
  - ▶ negotiate a session key (three messages are exchanged),
  - ▶ send the RMCP message protected with the sessions key over RSP.

## RSP session protocol (RSSP)



## The RMCP messages

- ▶ Presence Ping / Presence Pong;
- ▶ CapabilitiesRequest / CapabilitiesResponse;  
describe which operations are possible on the legacy and the secure ports;
- ▶ System State Request / System State Response;  
ask system about status, last boot medium etc.
- ▶ startup request (device can be specified: PXE, hard disk, CD-ROM);
- ▶ reboot request;
- ▶ stop request.

# Status of RMCP support

Some hardware with ASF support that we tested:

- ▶ HP Compaq dc7600
  - ▶ on the secure port, start/reboot/stop messages are processed;
- ▶ DELL Latitude D530 and Precision T5400
  - ▶ CapabilitiesRequest messages are processed,
  - ▶ CapabilitiesReply messages indicate that no remote administration function is supported.

## Remarks

- ▶ no vendor enabled remote administration on the legacy port;
- ▶ why do some vendors disable administration functions but still implement RMCP and SMBus functions?

## Card's behavior when using ASF

When it receives a packet, the card

- ▶ intercepts the packet;
- ▶ before transmitting it to the OS;
- ▶ checks if it is a RMCP packet;
- ▶ process it:
  - ▶ open/close session,
  - ▶ send system state informations,
  - ▶ perform system administration tasks,
  - ▶ the packet is **NOT** transmitted to the host.

# Protocol security (1/2)

## Potential issues

- ▶ protocol uses 160bit pre-shared keys, which means all clients might have the same keys;
- ▶ messages are integrity protected but the integrity pattern does not include message ID;
- ▶ in order to act as the *console*, an attacker just has to send a RAKP<sub>3</sub> with a valid HMAC.

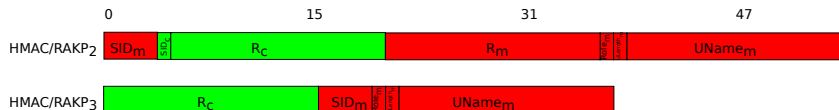
## Exploitation

- ▶ is it possible to forge the HMAC?
- ▶ can the client act as a integrity oracle?

## Protocol security (2/2)

### Not so easy

- ▶ only two concurrent sessions on the implementation we tested;
- ▶ not all fields are under the attacker's control;
- ▶ fields size problems.





# Implementation problems (1/2)

## Interesting fields under the attacker's control

- ▶ management console username (RAKP<sub>1</sub> message);
- ▶ management console session ID (Open Session Request).

## Trying to play with them, messing with

- ▶ size;
- ▶ content.

## Implementation problems (2/2)

### Username

- ▶ specifications limit the size to 16 chars, without NULL;
- ▶ the username size is coded using a 1 byte field.

### What if we don't play nice?

- ▶ "the card crashes";
- ▶ the host can only send Ethernet frames but not receive them.

What did exactly happen?

## Proof of concept

- ▶ Is the vulnerability serious?
- ▶ What are the direct and indirect consequences?
- ▶ How can we build a *proof of concept*?

Speaker switch

## Instrumenting the card (1/2)

### How did we find out what happened?

- ▶ what is really crashing?
- ▶ how is it crashing?

### Using the NetXtreme specifications

- ▶ public specs available for open-source developers;
- ▶ describe the internal card behavior;
- ▶ give informations about RX RISC components.

## Instrumenting the card (2/2)

How to debug the RX RISC? We need to:

- ▶ follow the execution flow;
- ▶ know registers contents;
- ▶ know why the CPU stops;
- ▶ trace data.

We need a *network card debugger*.

# What information can we retrieve from the card?

Specifications and experiments say we have access to:

- ▶ RX RISC mode register;
- ▶ RX RISC state register;
- ▶ RX RISC program counter;
- ▶ RX RISC hardware breakpoint register;
- ▶ some general registers.

We can use this knowledge to build that debugger.

## Our homemade debugger:

- ▶ uses information from mapped registers;
- ▶ runs in step-by-step mode;
- ▶ can perform register/memory tracking;
- ▶ can break on register/memory access;
- ▶ can do some pattern matching.

# Help

```
What should I do next (h for help)? h
```

```
Usage:
```

```
'a' -> Advance n steps  
's' -> Advance 1 step  
't' -> Trace  
'c' -> Continue  
'C' -> Continue (step-by-step)  
'g' -> Break on instruction  
'R' -> Break on pattern in register  
'S' -> Break on pattern in stack  
'H' -> Break on pattern in internal memory  
'M' -> Break on pattern in external memory  
'n' -> Break on next pattern in stack  
'l' -> Break on specific memory access  
'm' -> Break on any memory access  
'j' -> Break on register write  
'i' -> Break on instruction  
'T' -> Track register  
'L' -> Track memory address  
'Z' -> Track specific memory zone access  
'I' -> Track pattern in memory  
'P' -> Track pattern  
'x' -> Clear tracking  
'f' -> Find pattern in internal memory  
'F' -> Find pattern in external memory  
'A' -> Find all patterns in external memory  
'd' -> Display memory address  
'D' -> Display memory area  
'w' -> Write a word to memory address  
'r' -> Reset CPU  
'q' -> Quit
```



# CPU

```
***** Instruction
Instruction = 3c020001 LUI r2 = 00010000
Last memory access = 00000000
***** CPU Status Registers *****
RXPC      = 00011078 RXHWBRK = 0000001d
RXMODE    = 00009db0 RXSTATE = 80001400
```

# General registers

```
***** CPU Registers *****  
$0 = 00000000 $1 = 00010000 $2 = 00000000 $3 = 40000000  
$4 = 0001b4b8 $5 = 0001b8e6 $6 = 00000000 $7 = 0001bfc4  
$8 = 00000040 $9 = 00000050 $10 = 0001b8bc $11 = 0001bfc0  
$12 = 80000000 $13 = 00000001 $14 = 00000000 $15 = fffffbf  
$16 = a4020000 $17 = aaaaaaaaa $18 = 00000000 $19 = 0001af48  
$20 = 0000ad60 $21 = 018004f1 $22 = 000000fc $23 = 00010000  
$24 = ffffffff $25 = 80000000 $26 = 00000b50 $27 = 00011104  
$28 = c0000000 $29 = 0001bfd8 $30 = 0001c000 $31 = 000111f8
```

# Stack

```
***** Stack *****  
Stack pointer: 0001bfd8 (max) stack size: 10  
Stack bottom: 0001c000  
*****  
0001bffc:73fffffff 0001bfe8:00010e00  
0001bff8:00010044 0001bfe4:0001a918  
0001bff4:0001a000 0001bfe0:0000ad60  
0001bff0:0000ad64 0001bfdc:0001a000  
0001bfec:0001a80c 0001bfd8:00010b3c
```

## Why does the card crash?

RX RISC state register provides useful information:

1. bad memory alignment
2. invalid instruction fetch (jump to invalid location);
3. invalid data access (load/store in invalid location);
4. invalid instruction;

Points 2 and 4 can mean direct flow execution redirection.

Points 1 and 3 can mean indirect flow execution redirection (try to overwrite a return address in the stack).

## Changing the execution flow

When the RX RISC CPU is crashing, an attacker needs to:

- ▶ find the source of the data;
- ▶ tune it to fit her needs.

### Trials and errors

We managed to:

- ▶ make the *username* field overflow;
- ▶ overwrite a return address in the stack with an address under our control.

## Proof of concept code injection (1/2)

On this particular NIC and firmware version, an attacker is able to perform arbitrary code execution:

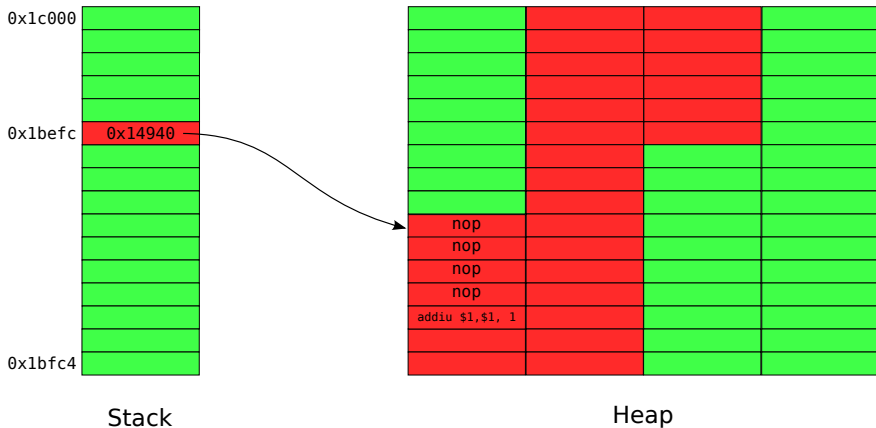
### Initial jump

- ▶ an attacker can overwrite a return address in the stack;
- ▶ she can find a stable (for a firmware version) memory address for *username*;
- ▶ she can put exploit code in *username* and jump there.

### Stage 1

- ▶ *username* is 255 chars (minus padding), not much instructions;
- ▶ but the attacker has access to network buffers;
- ▶ she can put code in a previously sent packet and jump there.

# Exploit and stage 1



## Proof of concept code injection (2/2)

### stage 2

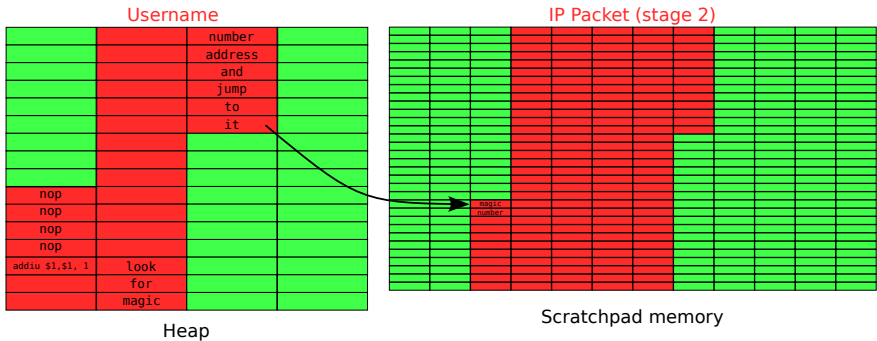
- ▶ size virtually unlimited;
- ▶ sent like a normal packet before the exploit;
- ▶ prepended by a magic number so stage 1 can find it.

### Now the attacker can:

- ▶ run arbitrary code on the RX RISC;
- ▶ provide new code using simple packets;
- ▶ rewrite the firmware if needed;
- ▶ ...



# Stage 2



## Man in the middle

Every packet ends up in the card memory:

- ▶ received packets before reaching the host;
- ▶ sent packets before being emitted on the wire.

### Play

- ▶ reroute DNS traffic;
- ▶ reroute *all* traffic;
- ▶ modify TLS negotiations;
- ▶ perform any conceivable MITM stealthily.

# Remote management

## Remember DELL disabled remote management?

- ▶ but the controller is connected to the SMBus;
- ▶ ASF! description table is present with the remote control functions;
- ▶ the exploit send messages to the SMBus;
- ▶ therefore it can perform *power-up*, *power-down*, *power-cycle*, ...

It can be reimplemented!

# Take-over the host

## The network card:

- ▶ is on the PCI/PCI-Express bus;
- ▶ can read/write to PCI configuration space;
- ▶ has Direct Memory Access (DMA) to the host.

The attacker taking over the NIC can read and write to main host memory!

# Using DMA

## DMA transfers

- ▶ NIC and host share network packets using DMA;
- ▶ meta-data (NIC address, host address, size) are stored in special structures, the *buffer descriptors*.

## proof of concept code: write to main memory

- ▶ write an host address to a buffer descriptor address field in the NIC;
- ▶ send packets;
- ▶ packet is written to the main memory at given address.

(almost) reliable.

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## OS dependent

- ▶ like all DMA-based attacks;
  - ▶ need to get around IOMMU;
  - ▶ need to find out where to read/write;
  - ▶ need to trigger the code execution.
- 
- ▶ for the proof of concept, we used Linux (because we know how it works);
  - ▶ same would work for any other OS;
  - ▶ nice trick, configure a new mac address on the NIC:  
90:90:90:90:90:90.

# Demonstration

## What we do in the demo:

- ▶ write some code at address 0 to run a remote shell;
- ▶ hook ourselves into `icmp_rcv` to jump at address 0;
- ▶ send a *magic ping*.

## Countermeasures

- ▶ use a patched firmware;
- ▶ deactivate ASF (not only in the BIOS);
- ▶ filter ASF and RMCP UDP ports;
- ▶ use an IOMMU on a supported OS;
- ▶ deactivate remote administration protocols, or
- ▶ reserve remote administration to safe/separated networks.
  - ▶ nobody ever enabled ASF on a laptop connected to Internet anyway
  - ▶ is it really safe to assume that?



## Conclusion

This vulnerability might seem scary, however remember:

- ▶ few cards support ASF;
- ▶ fewer cards enable ASF.

But,

- ▶ ASF is quite simple:
  - ▶ over UDP,
  - ▶ few cryptographic algorithms,
  - ▶ limited number of sessions,
  - ▶ no interaction with the network;
- ▶ AMT, IPMI, and the other remote management protocols are more complex:
  - ▶ over TCP,
  - ▶ heavy use of *webservices* (XML-RPC, SOAP, ...),
  - ▶ interactions with the whole network infrastructure (*Active Directory, Kerberos, ...*).

## Conclusion (2/2)

- ▶ more and more devices require firmwares:
  - ▶ network cards,
  - ▶ wireless network cards,
  - ▶ GSM and UMTS chipsets,
  - ▶ RAID controllers;
- ▶ with common characteristics:
  - ▶ no source code available,
  - ▶ close to the hardware,
  - ▶ possible access to the outside world (network cards),
  - ▶ real-time constraints.

More issues are likely to appear in the future.

It is time to develop simpler network cards and smaller drivers.

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## Question & answers

?

FAQ are available at

<http://www.ssi.gouv.fr/trustnetworkcard>