### **Blocking DNS Messages is Dangerous**

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- Created in 2009, the ANSSI is the French national authority for the defense and the security of information systems
  - in French, ANSSI, Agence Nationale de la Sécurité des Systèmes d'Information
- Under the authority of the Prime Minister
- Main missions are:
  - prevention
  - defense of French information systems
- One of its priorities is DDoS prevention and mitigation

### http://www.ssi.gouv.fr/en

## State of the art regarding DNS-related DDoS



Threat: on IP networks, sender address can be spoofed





### Principle:

- Based on reflection attacks
- Increase the attacker throughput by leveraging non-malicious nameservers
- DNS answer IP packets are often 40-50 times the size of the associated query IP packets
- 2 Mbps (attacker)  $\Rightarrow$  100 Mbps (target)



### What can an operator do?

DNS messages can be filtered at different levels:

L3 Drop packets

L4-7 Drop DNS datagrams or queries

L7 Response Rate Limiting (RRL):

- Identical DNS answers detection
- Bind, NSD, Knot
- ► Slips a truncated answer every X queries
  - e.g. 2 Mbps (attacker)  $\Rightarrow$  up to 2 Mbps (target)

# Can anti-DDoS technologies be useful for cache poisoning attacks?



### Cache poisoning attacks reminder

Principle:

Insert forged data in cache

Example:

2008: Kaminsky attack

Current Fix:

Source Port Randomization

Long Term Fix:

- DNSSEC
  - Requires large adoption





### Our cache poisoning attack: Step by step





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2: Perform the recursive resolution



### Our cache poisoning attack: Step by step



2 bis: Trigger anti-DDoS mechanism against the resolver



### Our cache poisoning attack: Step by step



3: Either answer with a truncated answer or drop the query **Dropping answers lead to resolver timeouts and retries** 



### Our cache poisoning attack: Step by step



3 bis: Send lots of Kaminsky-style answers to poison the cache

### Experiments & results



### A single authoritative nameserver

- Realistic thanks to authoritative nameserver selection attacks (Shulman fragment attacks, SRTT tricks...)
- A single outbound IP on resolver
- 100 Mbps of spoofed traffic
  - would go unnoticed by most ISP
- RRL with slip=2



### Validation of the theoretical model



#### Time

### We mathematically modeled the attack Details available on demand

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Based on the model, real-world attacks can be successful with a probability P in less than the following time estimates:

P: Probability of a successful cache poisoning attack

P = 10%	pprox 1h 15min
P = 50%	pprox 8h

### Are firewalls doing any better?







### **ANSSI** recommendations



### Always answer queries

# Never drop DNS queries when you can't tell which are legitimate

**Slip 1** is the only RRL safe configuration against our cache poisoning attack



### **Disclosure timeline**

#### **Timeline and feedbacks**

#### Disclosure timeline:

- June: DNS Software Vendors, Packagers
- August: NIC and root operators
- May-August: CERTs

### Security notifications:

- CVE-2013-5661 and CVE-2013-5752
- CERTA and NCSC advisory bulletin (September 9th, 2013)

### All have confirmed the vulnerability

### Some raised concerns

### Is slip 1 dangerous?





### courtesy of Pnetnod

As slip 1 grants an even payback, is this configuration dangerous for PPS attacks?





Facts:

- Current attacks are volumetric/bandwidth-related DDoS
- More susceptible protocols available for PPS attacks



# Network DDoS on the authoritative nameservers because of slip 1?



Facts:

- Amplification factor: 1:1
- Operators have symmetric bandwidth

### Investigation should be led if upload capacity is reached



## Computational DDoS on the authoritative nameservers because of slip 1?



#### Fact<sup>a</sup>:

 Slip 1 increases CPU consumption by less than 5% depending on implementations

atested on Xeon X5650 @2.67Ghz with 4000 qps





### courtesy of Pnetnod

# Network DDoS on the resolver because of slip 1 on the authoritative nameserver?





#### Fact:

- On average, the number of packets exchanged between a resolver and authoritative nameserver per query:
  - ► Slip 1: 9
  - ► Slip 2: 9.68





Computational DDoS on the resolver because of slip 1 on the authoritative nameserver?





Fact<sup>a</sup>:

 Slip 1 decreases CPU consumption by up to 20%, depending on implementations

atested on Xeon X5650 @2.67Ghz with 4000 qps



#### Summary

### RRL with Slip 1:

- Is worthless for attackers performing volumetric or PPS DDoS attacks
- ► is less CPU consuming for flooded resolvers
- Is a negligibly more CPU consuming for authoritative nameservers

### TL;DR summary: Slip 1 is OK



- Timeouts lead to more efficient cache poisoning attacks
- Always answering queries:
  - Thwarts our attack
  - Offers no benefit for attackers
- ► RRL Slip=1 mitigates DDoS
  - RRL Slip=2 is overkill for current DDoS attacks and is vulnerable to our cache poisoning attack
- Always answering is a temporary fix:
  - DNSSEC wake-up call?



### Thank you for your attention

Any questions?

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### Packets count for DOS of recursive servers

# $E(PC) = \sum_{i=1}^{n} \left(1 - \frac{1}{s}\right)^{i-1} \left(1 + \frac{8}{s}\right)$

with E(PC) being the mean packet count, n being the number of retries by the resolver and s being the slip value